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Graphs with the strong Havel–Hakimi property

Michael D. Barrus*

Grant Molnar†

Abstract

The Havel–Hakimi algorithm iteratively reduces the degree sequence of a graph to a list of zeroes. As shown by Favaron, Mahéo, and Saclé, the number of zeroes produced, known as the residue, is a lower bound on the independence number of the graph. We say that a graph has the strong Havel–Hakimi property if in each of its induced subgraphs, deleting any vertex of maximum degree reduces the degree sequence in the same way that the Havel–Hakimi algorithm does. We characterize graphs having this property (which include all threshold and matrogenic graphs) in terms of minimal forbidden induced subgraphs. We further show that for these graphs the residue equals the independence number, and a natural greedy algorithm always produces a maximum independent set.

1 Introduction

Given the degree sequence d of a simple graph, the *Havel–Hakimi algorithm* is the procedure that iteratively removes a largest term t from the sequence and subtracts 1 from the t largest remaining terms, stopping when a list of zeroes is produced. For example, starting with the degree sequence $d = (3, 2, 2, 2, 2, 1)$, the algorithm produces the sequences below.

$$d = d^0 : (3, 2, 2, 2, 2, 1)$$

$$d^1 : (2, 1, 1, 1, 1)$$

$$d^2 : (1, 1, 0, 0)$$

$$d^3 : (0, 0, 0)$$

Here d^i indicates the sequence obtained after i iterations of the removal/reduction step. Throughout the paper, we will order the terms of each d^i from largest to smallest.

This algorithm was introduced independently by Havel [8] and Hakimi [7] as a means for testing whether a list of nonnegative integers d is the degree sequence of a graph; d is a degree sequence if and only if the algorithm terminates at a list of zeroes. For a degree sequence d , the *residue* of d , denoted $R(d)$, is the number of zeroes that remain. For example, the steps above show that $R((3, 2, 2, 2, 2, 1)) = 3$. For any graph G having degree sequence d , we define $R(G) = R(d)$ and speak interchangeably of the residue of the graph and the residue of the degree sequence.

Let us say that a vertex v in a graph G has the *Havel–Hakimi property* if v is a vertex of maximum degree in G having neighbors with as large of degrees as possible; more precisely, v has the Havel–Hakimi property if it has maximum degree in G and no vertex adjacent to v has degree smaller than any vertex not adjacent to v . Each reduction from d^i to d^{i+1} in the Havel–Hakimi algorithm simulates the effect on the degree sequence of deleting a vertex with the Havel–Hakimi property.

Not every realization of a degree sequence has such a vertex. For example, there is no vertex with the Havel–Hakimi property in the realization of $(3, 2, 2, 2, 2, 1)$ obtained by attaching a pendant vertex to a vertex of a chordless 5-cycle. However, as shown in most proofs of the result of Havel and Hakimi (see, for

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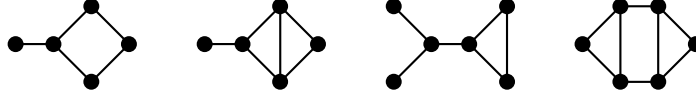


Figure 1: The 4-pan, kite, stool, and co-domino graphs.

instance, [2, Theorem 1.12]), given any degree sequence d and any vertex v chosen to have maximum degree, there is a realization of d where v has the Havel–Hakimi property.

We say that a graph G has the *strong Havel–Hakimi property* if in every induced subgraph H of G , every vertex of maximum degree has the Havel–Hakimi property. Let \mathcal{S} denote the class of graphs with this property. In this paper we study \mathcal{S} , characterizing its elements and describing some properties of these graphs.

First, we characterize \mathcal{S} in terms of its minimal forbidden induced subgraphs. Given a class \mathcal{F} of graphs, let $\text{Forb}(\mathcal{F})$ denote the class of all graphs having no induced subgraph isomorphic to an element of \mathcal{F} . Before presenting the set \mathcal{F} for which $\text{Forb}(\mathcal{F}) = \mathcal{S}$, we need some definitions.

Given graphs G and H , let $G + H$ denote the disjoint union of G and H , and express the disjoint union of m copies of G by mG . Given positive integers m and n , let P_n denote the path with n vertices, let K_n be the complete graph on n vertices, and let $K_{m,n}$ be the complete bipartite graph with partite sets of sizes m and n . Let $K_{2,3}^+$ be the graph obtained by adding an edge to the larger partite set of $K_{2,3}$; this graph is the complement of $K_2 + P_3$. We define the *4-pan*, *kite*, *stool*, and *co-domino*, respectively, to be the connected graphs shown in Figure 1.

Theorem 1. $\mathcal{S} = \text{Forb}(\{P_5, 4\text{-pan}, K_{2,3}, K_{2,3}^+, \text{kite}, 2P_3, P_3 + K_3, \text{stool}, \text{co-domino}\})$.

We prove this theorem in Section 2. Along the way, we see how \mathcal{S} compares with some other well known graph families. (A discussion of these families accompanies the proof.)

Corollary 2. *The class \mathcal{S} contains all matrogenic graphs (and hence all matroidal and threshold graphs as well).*

Further motivation for studying \mathcal{S} comes from the connection between the Havel–Hakimi residue and the independence number of a graph. It follows from the definition of the strong Havel–Hakimi property that when we iteratively delete vertices of maximum degree from a graph G in \mathcal{S} , we arrive at an edgeless induced subgraph of G with $R(G)$ vertices. Thus $R(G) \leq \alpha(G)$ for graphs G in \mathcal{S} , where $\alpha(G)$ denotes the independence number of G . Favaron et al. [4] showed that this inequality is in fact true for all graphs.

Theorem 3 ([4]; see also [6, 11]). *For every graph G , we have $R(G) \leq \alpha(G)$.*

This fact was first proposed by the conjecture-making computer program Graffiti due to Fajtlowicz [3], who also defined the residue as we use it here. Theorem 3 was generalized by Jelen in [9]; for a more detailed summary of these and related ideas, see [1].

Given that the residue of a graph G may be computed in $O(e)$ steps, where e is the number of edges in the graph, while determining $\alpha(G)$ is NP-hard, we may wish to know how well $R(G)$ approximates $\alpha(G)$. In particular, it is an open problem to find a nontrivial answer to the following question.

Question 1. *For which graphs are the residue and independence number equal?*

In considering when the residue of a graph equals its independence number, it seems natural to consider procedures that generate independent sets by removing high-degree vertices. The simplest such heuristic, called Maxine by Fajtlowicz (see [6]), begins with a graph G and, as long as the current graph contains any edges, iteratively chooses and deletes a vertex of maximum degree. The vertices that remain at the completion of the process necessarily form an independent set in G , and the size M of this set is clearly a lower bound on $\alpha(G)$. We note that different choices of a vertex of maximum degree can result in different sizes of independent sets; performing Maxine on the 5-vertex path can yield an independent set of size M equal to either 2 (which equals the residue) or 3 (which equals the independence number), depending on which vertex is deleted first. Still, Griggs and Kleitman proved the following:

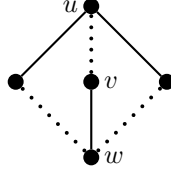


Figure 2: A configuration present when a vertex of maximum degree lacks the Havel–Hakimi property.

Theorem 4 ([6]). *If M is the size of the independent set produced by any application of the Maxine heuristic, then $R(G) \leq M \leq \alpha(G)$.*

Thus if the residue equals the independence number, the residue must equal the size of the independent set generated by any application of Maxine. As we now see, this is true for all graphs in \mathcal{S} .

Theorem 5. *If a graph G is in \mathcal{S} , then $R(G) = \alpha(G)$. Hence the Maxine heuristic always produces a maximum independent set for graphs in \mathcal{S} .*

A proof of this theorem is given in Section 3.

Throughout the paper, for a collection J of vertices in G , we use $G[J]$ to denote the induced subgraph of G with vertex set J , and we use $G - J$ to indicate the graph obtained by deleting all vertices in J from G .

2 Proofs of Theorem 1 and Corollary 2

This section contains a proof of Theorem 1, showing that the set

$$\mathcal{F}_{\mathcal{S}} = \{P_5, 4\text{-pan}, K_{2,3}, K_{2,3}^+, \text{kite}, 2P_3, P_3 + K_3, \text{stool}, \text{co-domino}\};$$

contains all minimal forbidden induced subgraphs for the class \mathcal{S} . As a byproduct, partway through the section we obtain a proof of Corollary 2, showing that \mathcal{S} contains all matrogenic graphs. We begin by showing that $\mathcal{S} \subseteq \text{Forb}(\mathcal{F}_{\mathcal{S}})$.

Lemma 2.1. *Every graph in \mathcal{S} is an element of $\text{Forb}(\mathcal{F}_{\mathcal{S}})$.*

Proof. Observe that in each graph in $\mathcal{F}_{\mathcal{S}}$, some vertex of maximum degree has a neighbor w and a non-neighbor u such that the degree of w is less than the degree of u . By definition, such vertices cannot occur in any induced subgraph of an element of \mathcal{S} ; thus graphs in \mathcal{S} must be $\mathcal{F}_{\mathcal{S}}$ -free. \square

We now show that $\text{Forb}(\mathcal{F}_{\mathcal{S}}) \subseteq \mathcal{S}$. Suppose to the contrary that $\text{Forb}(\mathcal{F}_{\mathcal{S}})$ contains an element G that is not in \mathcal{S} . By definition, some induced subgraph G' of G contains a vertex of maximum degree not having the Havel–Hakimi property. Since $\text{Forb}(\mathcal{F}_{\mathcal{S}})$ is a hereditary class and by definition $G' \notin \mathcal{S}$, it suffices to assume that $G' = G$. Some vertex v of maximum degree in G has a neighbor w and a nonneighbor u such that $d_G(w) < d_G(u)$. Because v is a neighbor of w that u does not share, u must be adjacent to two vertices to which w is not. We illustrate the adjacency relationships among u, v, w , and these other vertices in Figure 2, where dotted segments indicate non-adjacencies.

As a brief aside, we remark that the configuration illustrated in Figure 2 is precisely the forbidden configuration that characterizes the matrogenic graphs, as shown by Földes and Hammer in [5]. The matrogenic graphs were introduced in [5] as those graphs for which the vertex sets of a given 4-vertex configuration are the circuits of a matroid with ground set $V(G)$. The matrogenic graphs are known to properly contain such graph classes as the matroidal graphs and the threshold graphs (see [10] for a survey on all of these graph classes). No vertex deletion can create a copy of the configuration in Figure 2 where it was not already present; it follows that the class of matrogenic graphs is closed under taking induced subgraphs. We thus conclude that in every matrogenic graph each vertex of maximum degree has the Havel–Hakimi property. Hence \mathcal{S} contains the class of matrogenic graphs, as stated in Corollary 2.

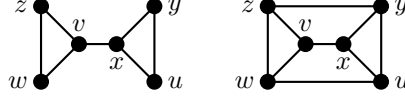


Figure 3: The graphs from Claim 3.

Continuing with our proof of Theorem 1, we use the assumption that G is \mathcal{F}_S -free to arrive at a contradiction in every case.

CLAIM 1: u and v have a common neighbor that is not adjacent to w .

Assume to the contrary that every common neighbor of u and v is adjacent to w . Let x and y be vertices adjacent to u and not to w . By assumption, neither x nor y is adjacent to v . Now since $d_G(v) \geq d_G(u) > d_G(w)$, v must have some neighbor z that is not adjacent to w . Note that z is distinct from u, w, x , and y . By assumption, z is not adjacent to u .

Suppose that x is not adjacent to y . Since $G[\{u, v, x, y, z\}]$ is not isomorphic to P_5 or the 4-pan, z is adjacent to neither x nor y . However, this would imply that $G[\{u, v, w, x, y, z\}]$ is isomorphic to $2P_3$ or that $G[\{u, v, w, x, z\}]$ is isomorphic to P_5 , depending on whether u is adjacent to w .

Suppose instead that x is adjacent to y . Now since $G[\{u, v, x, y, z\}]$ is not isomorphic to the kite graph, z is nonadjacent to at least one of x and y ; by symmetry, assume zx is not an edge of G . If uw is an edge, then $G[\{u, v, w, x, z\}]$ is isomorphic to P_5 , a contradiction, so u is not adjacent to w . If yz is an edge, then $G[\{v, w, x, y, z\}]$ is isomorphic to P_5 , again a contradiction, so y is not adjacent to z . However, now we see that $G[\{u, v, w, x, y, z\}]$ is isomorphic to the $P_3 + K_3$, a contradiction. Having arrived at a contradiction in every case, our proof of Claim 1 is complete.

We now establish a lemma that holds for all \mathcal{F}_S -free graphs.

CLAIM 2: Consider an \mathcal{F}_S -free graph H in which the subgraph induced on vertices a, b, c, d contains edges ab, bc, cd , and non-edges ac and bd . If some vertex p is adjacent to a but not to d , then p is adjacent to b and not to c .

Assume that p is adjacent to a but not to d . If p is not adjacent to b , then $H[\{a, b, c, d, p\}]$ is isomorphic to P_5 , 4-pan, or $K_{2,3}$. If p is adjacent to both b and c , then $H[\{a, b, c, d, p\}]$ is isomorphic to either kite or $K_{2,3}^+$, which proves Claim 2.

By Claim 1 above, G has a vertex x that is adjacent to u and v and not to w . Let

$$\begin{aligned} N_u &= \{p \in V(G) - \{u, v, w, x\} : pu \in E(G), pw \notin E(G)\}; \\ N_v &= \{p \in V(G) - \{u, v, w, x\} : pv \in E(G), px \notin E(G)\}; \\ N_w &= \{p \in V(G) - \{u, v, w, x\} : pw \in E(G), pu \notin E(G)\}; \\ N_x &= \{p \in V(G) - \{u, v, w, x\} : px \in E(G), pv \notin E(G)\}. \end{aligned}$$

Since G is \mathcal{F}_S -free, we may apply Claim 2 twice with $\{a, b, c, d\} = \{u, v, w, x\}$ to conclude that $N_w \subseteq N_v$ and $N_u \subseteq N_x$.

CLAIM 3: If $y \in N_u$ and $z \in N_v$, then $G[\{u, v, w, x, y, z\}]$ is one of the graphs shown in Figure 3.

Since $N_u \subseteq N_x$, vertex y is adjacent to x and not adjacent to v .

Suppose first that u is adjacent to neither w nor z . If y is adjacent to z , then since $G[\{v, w, x, y, z\}]$ is not isomorphic to 4-pan, we have that wz is an edge; however, then $G[\{u, v, w, x, y, z\}]$ is isomorphic to co-domino, a contradiction. Thus yz is not an edge. Now since $G[\{u, v, w, x, y, z\}]$ is not isomorphic to the stool graph, wz is an edge, and $G[\{u, v, w, x, y, z\}]$ is equal to the first graph shown in Figure 3.

Suppose instead that u is adjacent to z but not to w . Since $G[\{u, v, w, x, z\}]$ is not isomorphic to 4-pan, wz is an edge. Since $G[\{u, v, w, x, y, z\}]$ is not isomorphic to co-domino, yz is an edge. However, then $G[\{u, v, x, y, z\}]$ is isomorphic to $K_{2,3}^+$, a contradiction.

In light of the last two paragraphs, we may assume that uw is an edge. Since $G[\{u, v, w, x, z\}]$ is isomorphic to none of 4-pan, $K_{2,3}$, $K_{2,3}^+$, we see that uz is not an edge and wz is an edge in G . Now since $G[\{u, v, w, x, y, z\}]$ is not isomorphic to co-domino, yz must be an edge, and $G[\{u, v, w, x, y, z\}]$ is the second graph shown in Figure 3. Thus Claim 3 is true.

Inspecting the neighbors of z in both graphs in Figure 3, we conclude that if N_u is nonempty, then $N_v = N_w$. Now since $d_G(u) > d_G(w)$, certainly N_u is nonempty. However, recalling our assumptions on the degrees of u, v , and w , and that $N_u \subseteq N_x$, we have

$$|N_w| < |N_u| \leq |N_x| \leq |N_v| = |N_w|.$$

This contradiction implies that $\text{Forb}(\mathcal{F}_S) \subseteq \mathcal{S}$, which completes our proof of Theorem 1.

3 Proof of Theorem 5

In this section we prove Theorem 5 from the Introduction. We preface our main argument with results on the Maxine heuristic. First is the observation made there about Maxine and the residue, reproduced here for easy reference.

Observation 3.1. *If G is a graph and at each step of an application of the Maxine heuristic to G the maximum-degree vertex to be deleted has the Havel–Hakimi property, then the independent set returned by Maxine has size $R(G)$. Consequently, this conclusion holds for all graphs in \mathcal{S} .*

Lemma 3.2. *Let G be a graph having at least one edge. If a vertex v of maximum degree in G belongs to every maximum independent set in G , then v either belongs to an induced 4-cycle in G , or v is the center vertex of an induced P_5 .*

Proof. Suppose that v is a vertex of maximum degree in G and that v belongs to every maximum independent set in G . Let I be a maximum independent set in G , let $A = I - \{v\}$, and let $B = V(G) - I$. Since G has an edge and v is a vertex of maximum degree, v has at least one neighbor in G ; note that all neighbors of v belong to B . Since v belongs to every maximum independent set, using one of these neighbors to replace v in the set I cannot yield an independent set, so each neighbor of v has at least one neighbor in A . Since v also has maximum degree in G and all its neighbors have degree at least 2, v has degree at least 2 as well.

If the neighbors of v are all pairwise adjacent, then each has a degree larger than that of v , a contradiction. Thus v has two neighbors w and x that are nonadjacent. If w and x have a common neighbor in A , then this vertex and v, w, x induce a 4-cycle in G . If w and x have no common neighbor, then, since each has a neighbor in A , these neighbors and v, w, x induce a copy of P_5 in G with v as the center vertex. \square

In what follows we use C_4 to indicate a cycle on four vertices.

Corollary 3.3. *The Maxine heuristic always produces a maximum independent set when applied to a $\{C_4, P_5\}$ -free graph.*

Proof. It follows from Lemma 3.2 that deleting a vertex of maximum degree in a $\{C_4, P_5\}$ -free graph G cannot reduce the independence number of G . Inductively applying this result (since the resulting subgraphs are also $\{C_4, P_5\}$ -free), we conclude that the independent set resulting from the Maxine algorithm has size $\alpha(G)$. \square

We now complete our proof of Theorem 5. In the following, let G be an element of \mathcal{S} . We recall from Theorem 1 that G is $\{P_5, 4\text{-pan}, K_{2,3}, K_{2,3}^+, \text{kite}, 2P_3, P_3 + K_3, \text{stool}, \text{co-domino}\}$ -free, and we will refer to these forbidden subgraphs often in what follows.

We proceed by induction on the number of vertices in G . One easily verifies that $R(G) = \alpha(G)$ for all graphs on four or fewer vertices, and we suppose that G has n vertices and that $R(H) = \alpha(H)$ for all graphs H in \mathcal{S} with fewer than n vertices.

Let v be a vertex of maximum degree in G . Since G belongs to \mathcal{S} , the vertex v has the Havel–Hakimi property, and $G - v$ belongs to \mathcal{S} . Using the induction hypothesis, we have $R(G) = R(G - v) = \alpha(G - v)$. If $\alpha(G - v) = \alpha(G)$ our proof is complete, so suppose henceforward that v belongs to every maximum independent set in G . Since G is P_5 -free, Lemma 3.2 implies that v belongs to an induced 4-cycle in G . Let W be the vertex set of this cycle, and let U be the set of all vertices of G that are adjacent to every vertex of W .

Since G is $\{K_{2,3}, K_{2,3}^+\}$ -free, $G[U]$ is $\{3K_1, K_2 + K_1\}$ -free. It is an easy exercise to show that a graph is complete multipartite if and only if it is $\{K_2 + K_1\}$ -free. Thus $G[U]$, and by extension $G[U \cup W]$ is a complete multipartite graph where the maximum size of a partite set is 2. Let B be the union of all partite sets of size 1, and let C be the union of all partite sets of size 2 (so $W \subseteq C$). Finally, let A be the subset of $V(G) - B - C$ consisting of all vertices not having any neighbors in C , and let $D = V(G) - A - B - C$.

Any four vertices comprising two partite sets in $G[C]$ induce C_4 , and since G is $\{4\text{-pan}, K_{2,3}, K_{2,3}^+\}$ -free, each vertex in D is adjacent to either no vertex, all vertices, or exactly two consecutive vertices of such a 4-cycle. By definition each vertex in D has both a neighbor and a nonneighbor in W (otherwise the vertex would belong to A or $B \cup C$); it follows that vertices in D each have exactly one neighbor in each partite set of $G[C]$.

Since G does not induce P_5 in $G[W \cup \{a, d\}]$ for any vertices $a \in A$ and $d \in D$, no edges exist joining vertices in A with vertices in D . Likewise, since G does not induce the kite graph in $G[W \cup \{b, d\}]$ for any vertices $b \in B$ and $d \in D$, all edges possible between B and D must exist in G . Now any vertex in B has degree at least $|B| + |C| + |D| - 1$, while any vertex of C has degree at most $|B| + |C| + |D| - 2$. Since v belongs to C and is a vertex of maximum degree in G , we see that $B = \emptyset$.

Let u be the other vertex in the partite set of $G[C]$ that contains v . Every maximum independent set in G must contain a vertex from D that is adjacent to u ; otherwise, we could replace v by u in the set and obtain a maximum independent set in G that does not contain v , a contradiction. Thus D is nonempty, and we conclude that $G[C]$ is isomorphic to C_4 , since if a vertex d of D were adjacent to exactly one vertex of each of three different partite sets in $G[C]$, then the induced subgraph with vertex set consisting of d , one of the partite sets, and the non-neighbors of d in the other two partite sets is isomorphic to the kite, a contradiction.

Partition D into two sets D_u, D_v according to which of u and v each vertex is adjacent to. As observed above, D_u is nonempty. Since v has maximum degree in G , we see that $|D_v| \geq |D_u| \geq 1$. Let $x \in D_u$ and $y \in D_v$.

Now let $\{s, t\}$ be the partite set different from $\{u, v\}$ in $G[C]$. Each of x and y is adjacent to exactly one of s and t . Without loss of generality, assume that x is adjacent to s .

If y is also adjacent to s , then since $G[\{t, u, v, x, y\}]$ is not isomorphic to P_5 , vertices x and y must be adjacent. However, then $G[\{s, t, u, x, y\}]$ is isomorphic to the kite graph, a contradiction. Thus y is adjacent to t . By similar arguments, all vertices of D_u are adjacent to s , and all vertices of D_v are adjacent to t .

Each vertex in D_u must be adjacent to every vertex of D_v , since G is co-domino-free. Furthermore, D contains no vertex other than x or y , since a third vertex, together with x, y, u , and v , would form the vertex set of an induced 4-pan or kite. Thus $G[C \cup D]$ is isomorphic to the graph on the right in Figure 3, which contradicts the claim that v belongs to every maximum independent set in G , since a maximum independent set in G will include precisely two vertices from $G[C \cup D]$, and any two nonadjacent vertices will do, as there are no edges between A and D . We conclude that v cannot belong to every maximum independent set, and in fact $R(G) = R(G - v) = \alpha(G - v) = \alpha(G)$, which completes our induction.

Acknowledgments

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